

# MAS114: Lecture 6

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2020–2021

# Some news on assessment

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It's been decided that, to pass the module, you *must* pass both of these components.

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- then the statement  $P(n)$  is true for all  $n \in \mathbb{N}$ .



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When we are trying to prove the induction step  $P(k) \Rightarrow P(k + 1)$  we refer to  $P(k)$  as the *induction hypothesis*.

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We'll prove many things by induction in this course, but this is one:

### Proposition

*For any natural number  $n$ , we have the following formula for the sum of the first  $n$  positive integers:*

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Notice that  $P(n)$  is *not a number*, it's a *statement*.

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For our base case,  $P(0)$  says that the sum of *no integers at all* is  $0 \times 1/2$ , which is true, as the sum of no integers is zero.

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The statement  $P(k)$  tells us that

$$1 + 2 + \cdots + (k - 1) + k = \frac{k(k + 1)}{2}.$$

We need to prove  $P(k + 1)$ , which would say that

$$1 + 2 + \cdots + (k - 1) + k + (k + 1) = \frac{(k + 1)(k + 2)}{2}.$$

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Now note that

$$\begin{aligned} & 1 + 2 + \cdots + (k - 1) + k + (k + 1) \\ = & (1 + 2 + \cdots + (k - 1) + k) + (k + 1) \\ = & \frac{k(k+1)}{2} + (k + 1) \quad (\text{by the induction hypothesis}) \\ = & \frac{k(k+1)+2(k+1)}{2} \\ = & \frac{(k+1)(k+2)}{2}. \end{aligned}$$

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This is exactly the statement  $P(k + 1)$ , which is what we needed for the induction step, and that completes the proof. □

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If you believe that induction is a reliable method of proof (and I do, and I hope you do too), then it had better be the case that we're not using induction correctly.

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If you don't have a base case, such as  $P(0)$ , then it's of no use to prove that  $P(k) \Rightarrow P(k + 1)$  for all  $k$ . It's no use to be able to climb a ladder if the bottom of the ladder is unreachable.



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We'll take  $P(1)$  as the base case of the induction. This is the statement "Given any one horse, all of them have the same colour": this is obviously true.



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Now we'll prove the induction step. We will assume that  $P(k)$  is true ("given any  $k$  horses, all of them have the same colour"): our job is to prove that  $P(k + 1)$  is true ("given any  $(k + 1)$  horses, all of them have the same colour").

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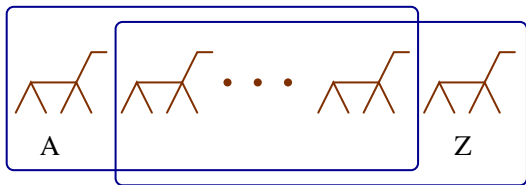
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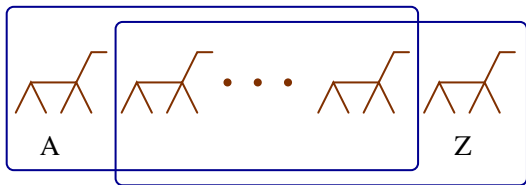
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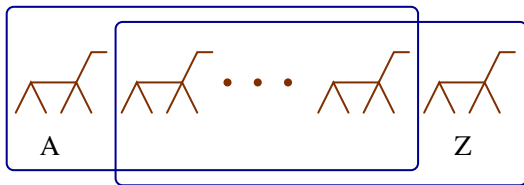


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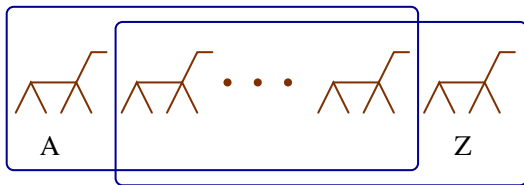


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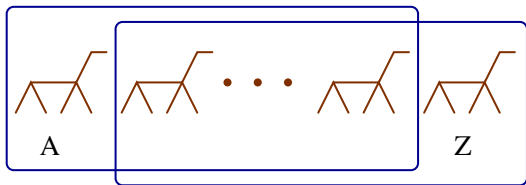
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Now we'll prove the induction step. We will assume that  $P(k)$  is true ("given any  $k$  horses, all of them have the same colour"): our job is to prove that  $P(k + 1)$  is true ("given any  $(k + 1)$  horses, all of them have the same colour").

So suppose we have  $(k + 1)$  horses. Name two of them Alice and Zebedee.



Excluding Alice, there are  $k$  horses, which all have the same colour, by the induction hypothesis. So all the horses except Alice have the same colour as Zebedee.

Also, excluding Zebedee, there are  $k$  horses, which all have the same colour, again by the induction hypothesis. So all the horses except Zebedee have the same colour as Alice.



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In fact, it's a parody of a *valid* style of argument. If it is the case that *any two things are the same*, then we could prove using exactly this method that they're *all the same*. In fact, this is something you already know, since “all are alike” and “no two differ” are synonymous phrases.



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(ii)  $P(k)$  implies  $P(k + 1)$  for all  $k \geq 15$ ,

then  $P(n)$  is true for all  $n \geq 15$ .

Perhaps you want to think of that as saying “if have a door which leads to the fifteenth rung of a ladder, and you know how to climb ladders, then you can get to every rung above the fifteenth”.

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Actually, you should have been prepared for this variant: my induction proof that “all horses have the same colour” started with 1, not 0. (Okay, that proof was wrong. But there was nothing wrong with *that bit* of the proof: there’s nothing wrong with induction starting from 1. It was something else that was wrong).